

Chapter 27 Modular Arithmetic

Def Let $m, n \in \mathbb{Z}$ be integers not both 0. Then their greatest common divisor is an integer $d \in \mathbb{Z}^+$ such that (1) $d|m$ and $d|n$ (common divisor)
(2) if $s \in \mathbb{Z}^+$ and $s|m$ and $s|n$, then $s|d$.
(greatest common divisor)

We denote d by $\gcd(m, n)$. We say m and n are relatively prime if $\gcd(m, n) = 1$.

Example Find $\gcd(-16, 40)$, $\gcd(0, 45)$, $\gcd(-30, -27)$.
 $= 8$ $= 45$ $= 3$

Theorem Let $m, n \in \mathbb{Z}$ be integers not both 0. Then their greatest common divisor exists, is unique, and $\gcd(m, n) = km + ln$ for some $k, l \in \mathbb{Z}$.

Proof Let $A = \{ xm + yn : x, y \in \mathbb{Z} \text{ and } xm + yn > 0\}$.

We claim A is nonempty. Indeed, $m^2 + n^2 \in A$ since $m^2 + n^2 > 0$ and $m^2 + n^2 = xm + ny$ when $x=m$ and $y=n$. Since $A \subseteq \mathbb{N}$ is non-empty, by the Well-Ordering Principle of \mathbb{N} , it has a minimum.

Let $d = \min A$. Since $d \in A$, $d = km + ln$ for some

$k, l \in \mathbb{Z}$. We claim $d = \gcd(m, n)$. First note $d \in \mathbb{Z}^+$

since $d \in A$. Next we must show $d|m$ and $d|n$.

By the division algorithm, $m = qd + r$ for some

$q, r \in \mathbb{Z}$ such that $0 \leq r < d$. We claim $r=0$.

Suppose not. Then $r > 0$ and

$$\begin{aligned} r &= m - qd \\ &= m - q(km + ln) \\ &= (1-qk)m + (-ql)n, \end{aligned}$$

which implies $r \in A$. However since $r < d$, this

contradicts the fact that $d = \min A$.

So since $r=0$, $m = qd$, which means $d|m$.

The same argument shows $d|n$. Next we must

show that if $s \in \mathbb{Z}^+$ such that $s|m$ and $s|n$,

then $s|d$. Exercise Prove that if $a|b$ and $a|c$

then $a|(bx+cy)$ for any $x, y \in \mathbb{Z}$ and use this here.

Finally, we must show uniqueness: If $t \in \mathbb{Z}^+$ satisfies properties

(1) and (2) of the gcd definition, then $t=d$.

Exercise Prove that if $a, b \in \mathbb{Z}^+$ such that $a|b$

and $b|a$, then $a=b$ and use this here.

Euclidean Algorithm Let $m, n \in \mathbb{Z}^+$ and suppose $m > n$.

Repeatedly apply the division algorithm:

$$\begin{aligned} m &= q_1 n + r_1 \\ n &= q_2 r_1 + r_2 \\ r_1 &= q_3 r_2 + r_3 \\ r_2 &= q_4 r_3 + r_4 \\ &\vdots \\ r_{k-3} &= q_{k-1} r_{k-2} + r_{k-1} \\ r_{k-2} &= q_k r_{k-1} + r_k \\ r_{k-1} &= q_{k+1} r_k + 0. \end{aligned}$$

Eventually you'll get a remainder of 0 since

$n > r_1 > r_2 > \dots \geq 0$. Then $r_k = \gcd(m, n)$.

Example Use the Euclidean Algorithm to compute

$\gcd(35, 20)$ and $\gcd(86, 24)$.

$$\begin{array}{l} 35 = 1 \cdot 20 + 15 \\ 20 = 1 \cdot 15 + 5 \\ 15 = 3 \cdot 5 + 0 \end{array} \quad \left\{ \Rightarrow \gcd(35, 20) = 5 \right. \quad \begin{array}{l} 86 = 3 \cdot 24 + 14 \\ 24 = 1 \cdot 14 + 10 \\ 14 = 1 \cdot 10 + 4 \\ 10 = 2 \cdot 4 + 2 \\ 4 = 2 \cdot 2 + 0 \end{array} \quad \left. \right\} \quad \begin{array}{l} \gcd(86, 24) = 2. \end{array}$$

Why the algorithm works:

(1) $r_k \mid r_{k-1}$ since $r_{k-1} = q_{k+1} r_k$

$\Rightarrow r_k \mid r_{k-2}$ since r_{k-2} is a linear combination
of r_{k-1} and r_k and
 $r_k \mid r_{k-1}$ and $r_k \mid r_k$.

$\dots \Rightarrow r_k \mid m$ and $r_k \mid n$.

So $r_k \mid \gcd(m, n)$.

(2) Let $d = \gcd(m, n)$. Then $d \mid m$ and $d \mid n$

and $r_1 = m - q_1 n \Rightarrow d \mid r_1$

Since $d \mid n$ and $d \mid r_1$ and $r_2 = n - q_2 r_1$

$\Rightarrow d \mid r_2$.

$\dots \Rightarrow d \mid r_k$.

So $d \mid r_k$ and $r_k \mid d \Rightarrow r_k = d$.